

Query languages for property graphs: RPQs in theory and practice

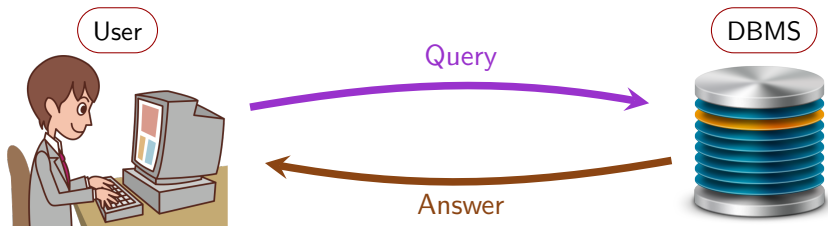
Victor MARSAULT

Université Gustave-Eiffel, CNRS, LIGM

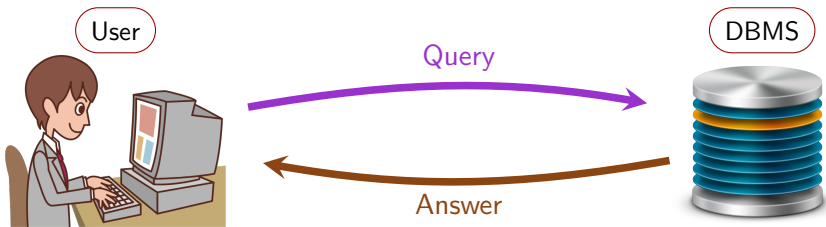
BAAM/ADA/MOA Seminar

2022-11-29

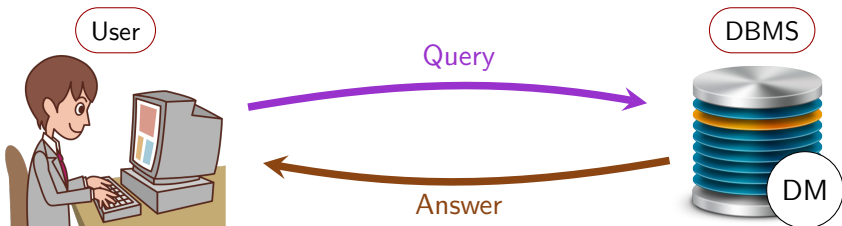
Introduction



- DBMS = **D**ata**B**ase **M**anagement **S**ystem

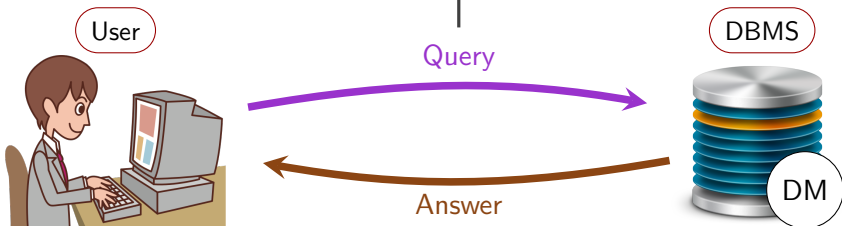


- DBMS = **D**ata**B**ase **M**anagement **S**ystem



- DM = **D**ata **M**odel = "*The way data is structured*"
 - Relational ? XML ? Property graph ? RDF ? etc.

- DBMS = **D**ata**B**ase **M**anagement **S**ystem
- Query language
 - "What can users ask for?"

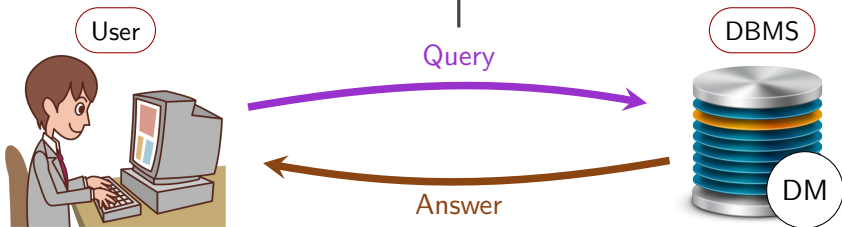


- DM = **D**ata **M**odel = "The way data is structured"
 - Relational ? XML ? Property graph ? RDF ? etc.

- DBMS = **D**ata**B**ase **M**anagement **S**ystem

- Query language

- *"What can users ask for?"*



- Semantics of query

- *"What does the query mean?"*

- DM = **D**ata **M**odel = *"The way data is structured"*

- Relational ? XML ? Property graph ? RDF ? etc.

Example: DB for a small store

Client table

name	address
Alice	Wonderland
Bob	124 Conch St.
Charlie	1593 Broadway

Product table

name	price
Sponge	1€
Broom	5€
Rabbit	0€
Pocket Watch	100€

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Order table

id	buyer	date
0	Alice	01-11-1865
1	Bob	07-07-2022

Order-content table

order_id	product
0	Rabbit
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Relational DBMS = tables with cross-references



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Vast majority of DMBS's are relational, not graph

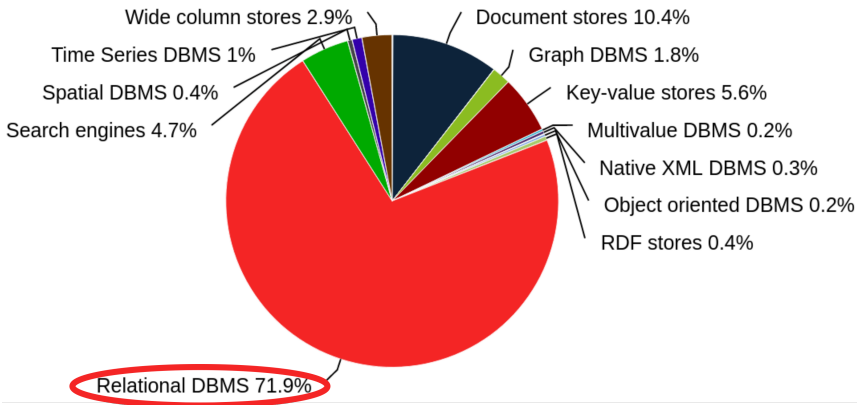


Figure and data from db-engines.com, June 2022

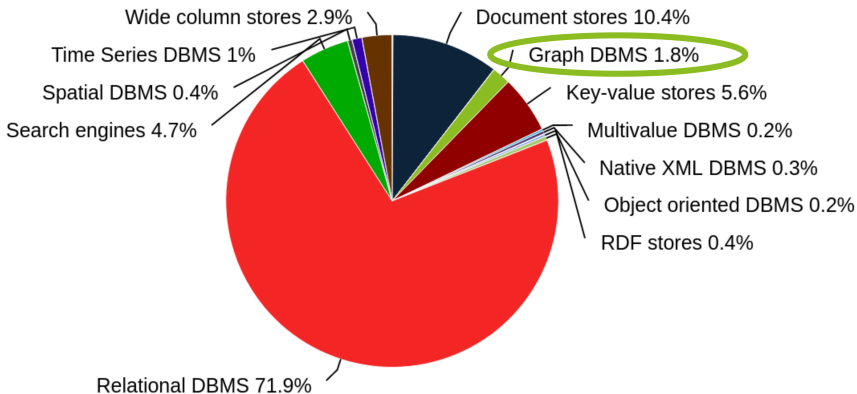


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Graph DBMS is growing in popularity



+25% per year since 2013

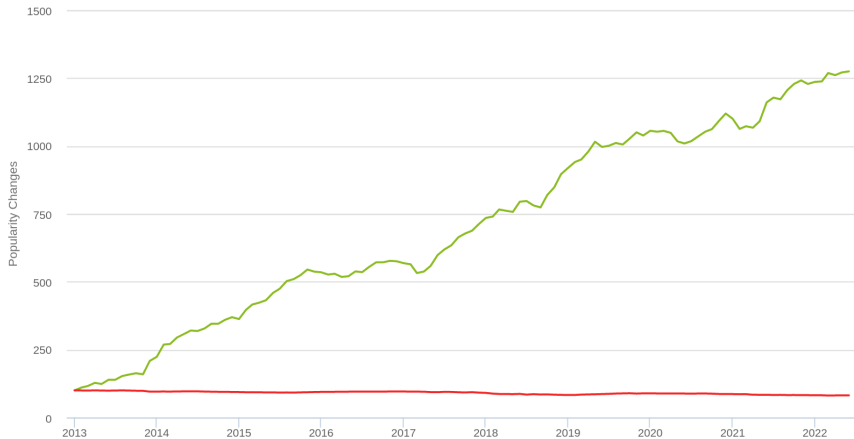
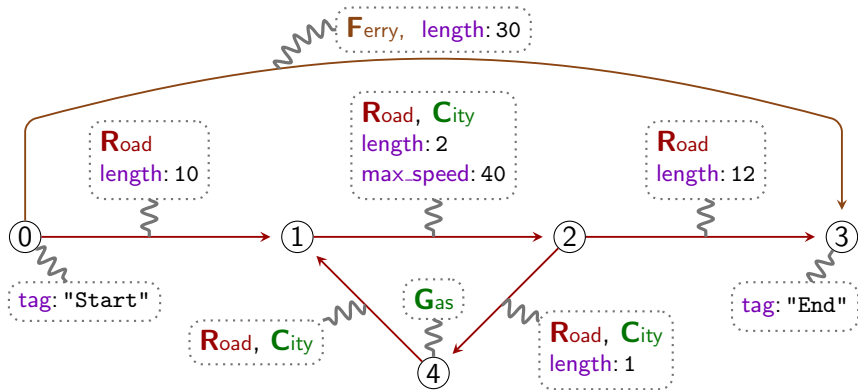


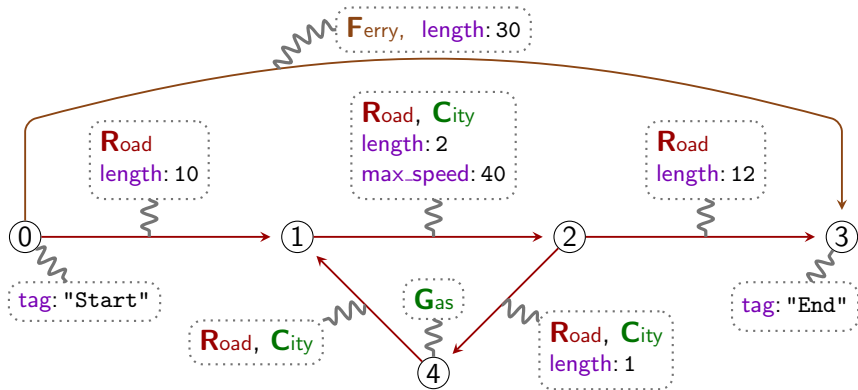
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Why use graph databases ?

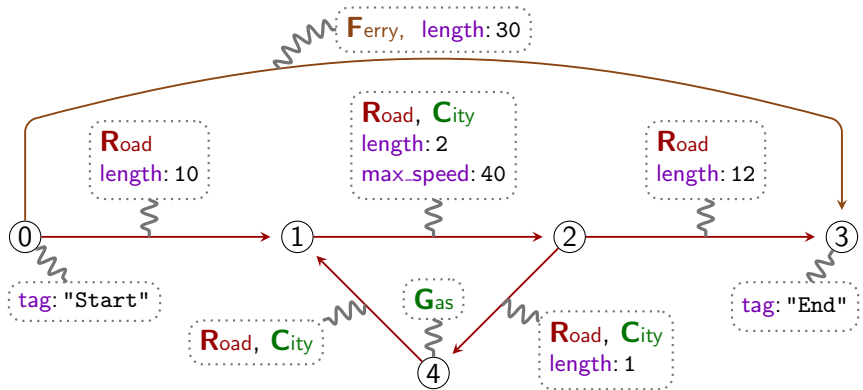
Some data have intrinsically the structure of graphs (e.g. networks)



Why not store graphs in tables?



Why not store graphs in tables?



id	source_id	target_id	Road	Ferry	City	length	max_speed
e ₀₁	0	1	true	false	false	10	
e ₁₂	1	2	true	false	true	10	40
e ₄₁	4	1	true	false	true		
⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮

GraphLabelsProperties

History of query languages for property graphs

1987 – RPQs are invented [Cruz-Mendelzon-Wood, 1987]



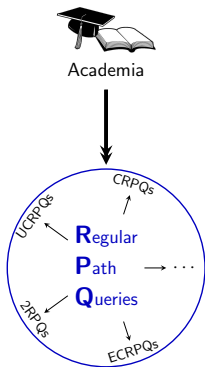
Academia



Regular
Path
Queries

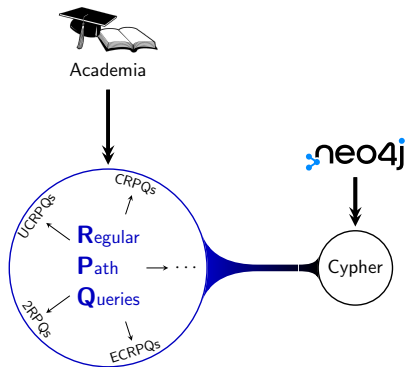
From RPQs to GQL: history and actors

Since 1990's – RPQs are extended and studied in academia



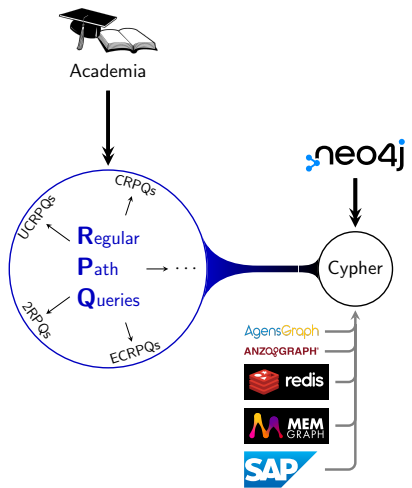
From RPQs to GQL: history and actors

2011 – Cypher is designed by Neo4j



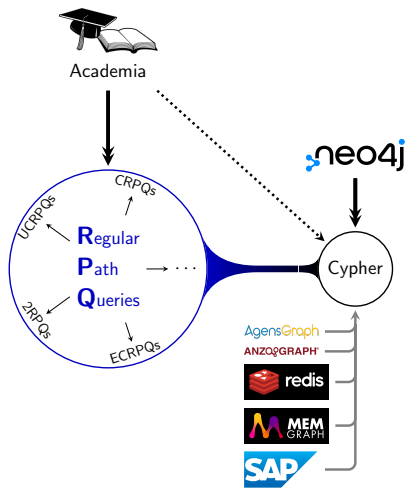
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mid 2010's – Cypher is becoming a standard de facto. Standardize Cypher?



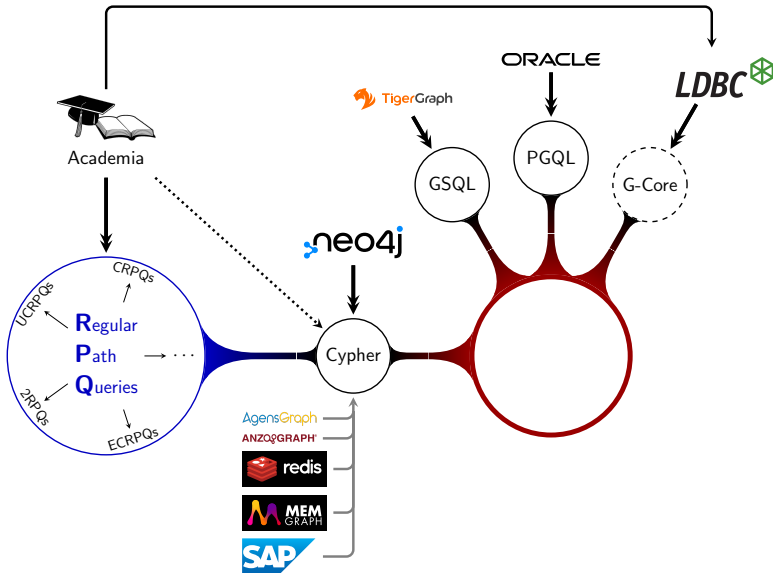
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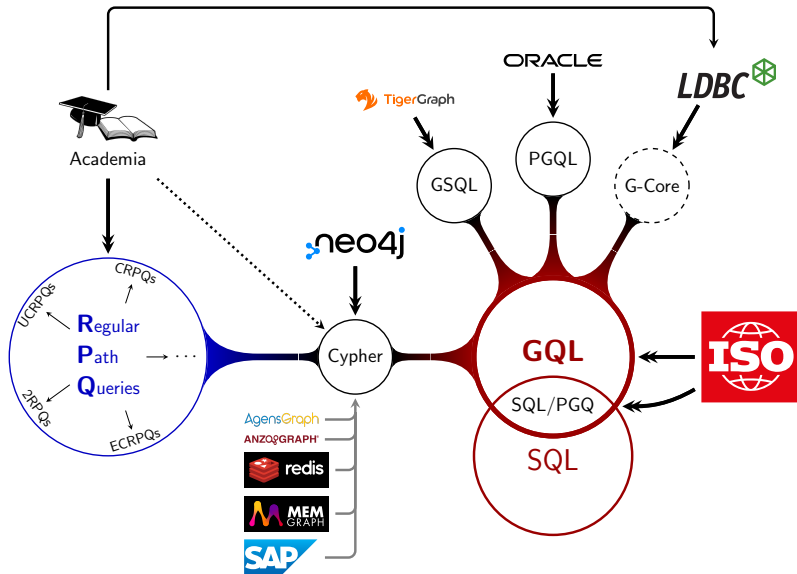
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late 2010's – Merge all existing languages instead of standardizing Cypher?



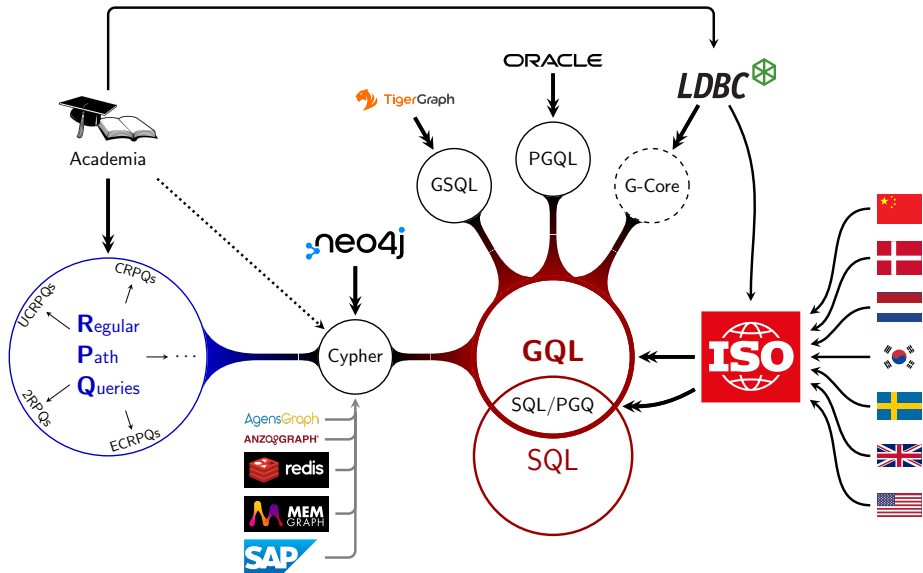
From RPQs to GQL: history and actors

2019-2021 – Two ISO projects: GQL [39075] and SQL/PGQ [9075-16.2]



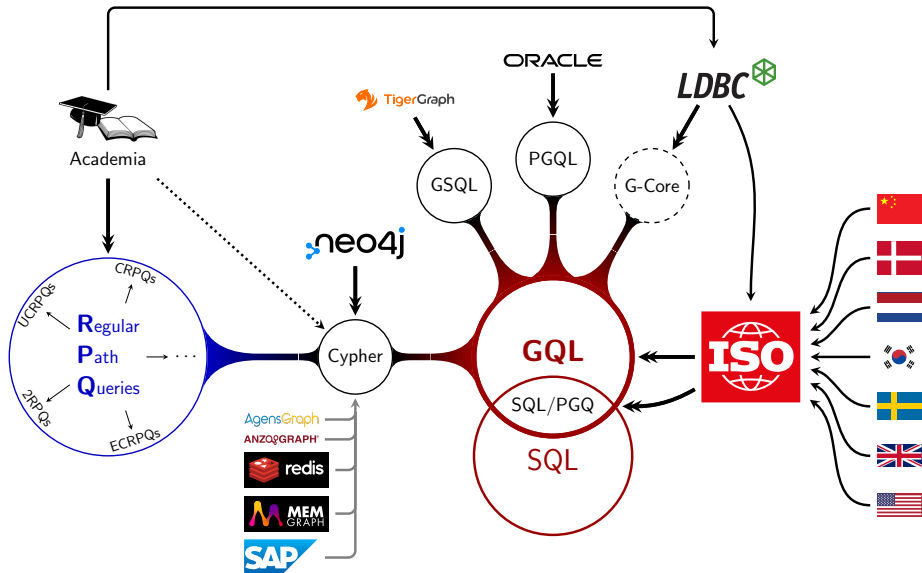
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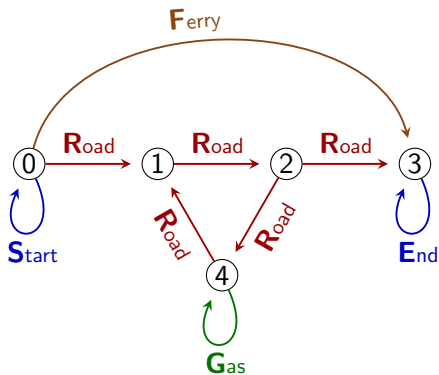
From RPQs to GQL: history and actors

2023 (expected) – Publication of version 1 of GQL



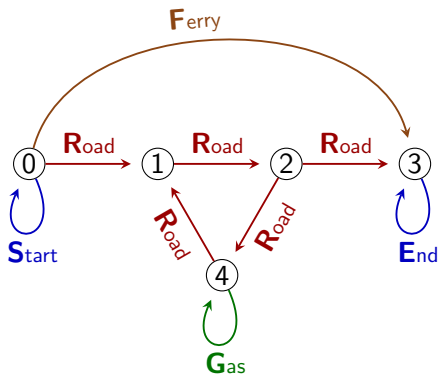
Foundation of querying graph databases: RPQs

A graph consists of ...



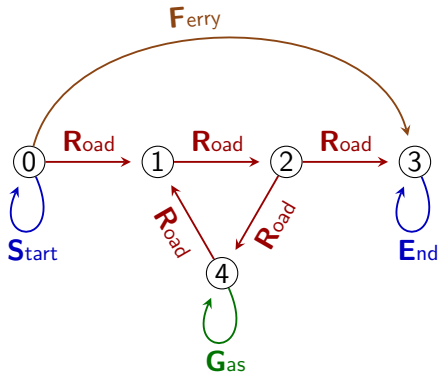
A graph consists of ...

- Vertices (or Nodes)



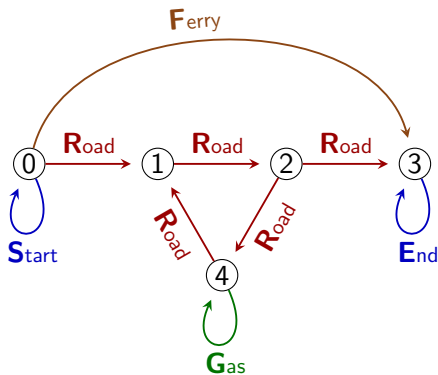
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A graph consists of ...

- Vertices (or Nodes)
- Edges (or Relationships)
- Edge labels: {**R**, **F**, **G**, **S**, **E**}

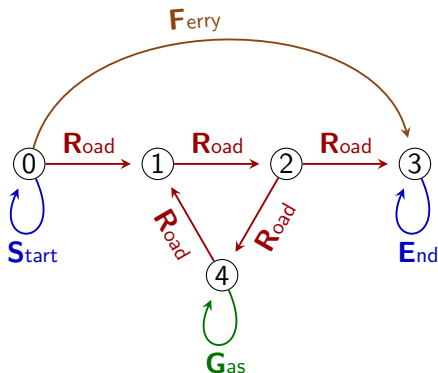


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Walk

- a.k.a. Path
- Sequence of edges
- Can reuse vertices and edges



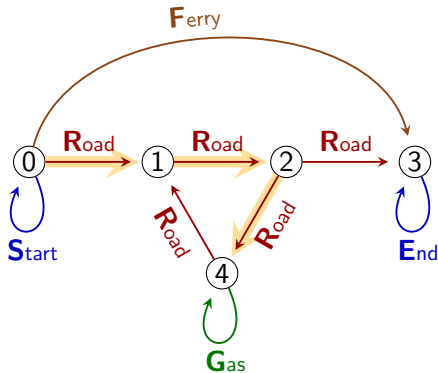
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0 → 1 → 2 → 4



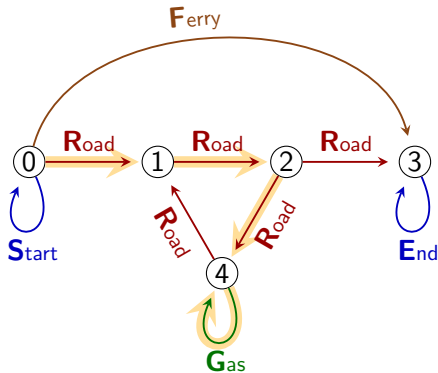
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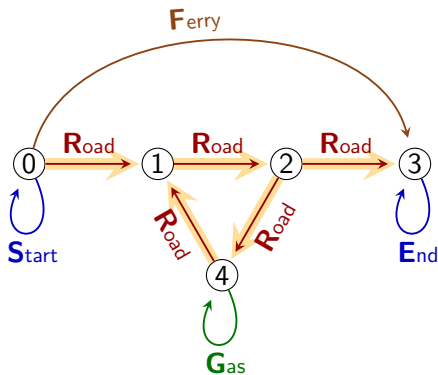
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Walk

- a.k.a. Path
- Sequence of edges
- Can reuse vertices and edges

0 → 1 → 2 → 4 → 1 → 2 → 3



$$Q ::= \mathbf{A}$$
$$QQ$$
$$Q + Q$$
$$Q^*$$

where \mathbf{A} is a label in the graph.

An RPQ denotes a set of words

$\mathbf{S}(\mathbf{F} + \mathbf{R})^* \mathbf{E}$ denotes the words of the shape $\mathbf{S} \underbrace{\langle \text{something} \rangle}_{\text{Any number of } \mathbf{F} \text{ and } \mathbf{R}, \text{ in any order}} \mathbf{E}$

$$Q ::= A$$

$$QQ$$

$$Q + Q$$

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An RPQ denotes a set of words

S (**F** + **R**)^{*} **E** denotes the words of the shape **S** <something> **E**

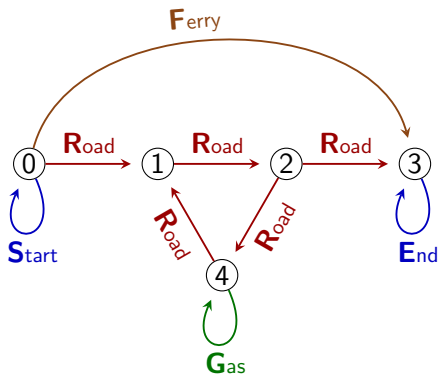
Any number of **F** and **R**, in any order

Matches

A match for Q is any walk w such that Q denotes the label of w

Query **RR** matches...

...walks of two **R_{oad}**-edges



Query **RR** matches...

...walks of two **R_{oad}**-edges

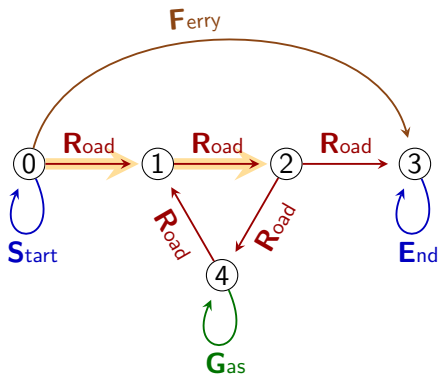
0 → 1 → 2

1 → 2 → 3

1 → 2 → 3

2 → 4 → 1

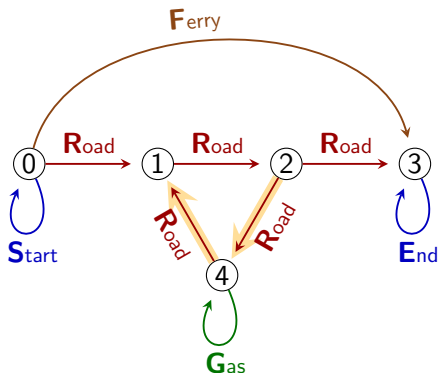
4 → 1 → 2



Query **RR** matches...

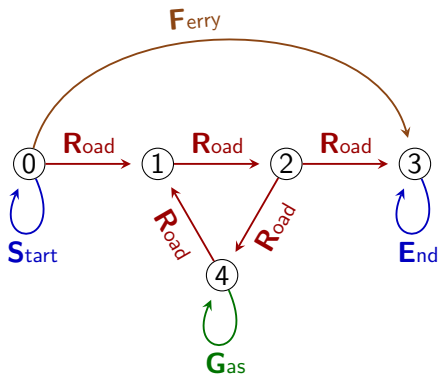
...walks of two **R_{road}**-edges

0 → 1 → 2	2 → 4 → 1
1 → 2 → 3	4 → 1 → 2
1 → 2 → 3	



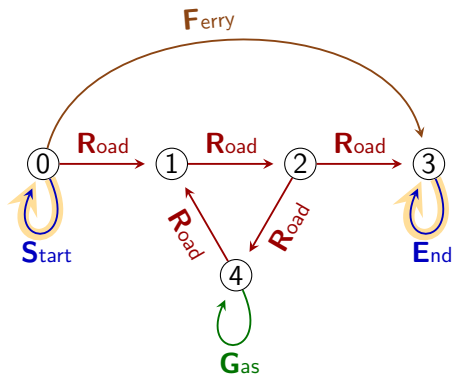
$$Q_1 = S(R+F)^*E$$

Q_1 matches...



$$Q_1 = \mathbf{S(R+F)^*E}$$

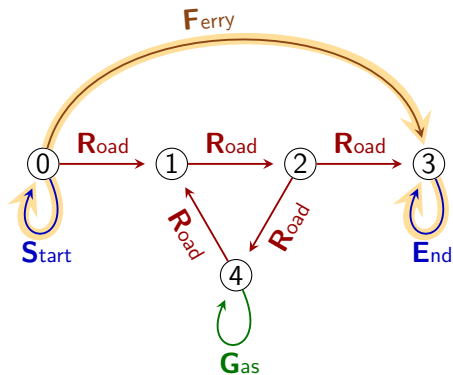
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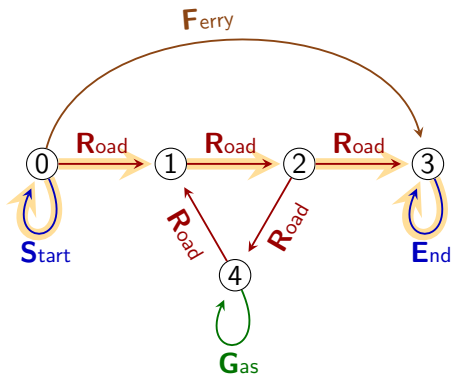
- The ferry



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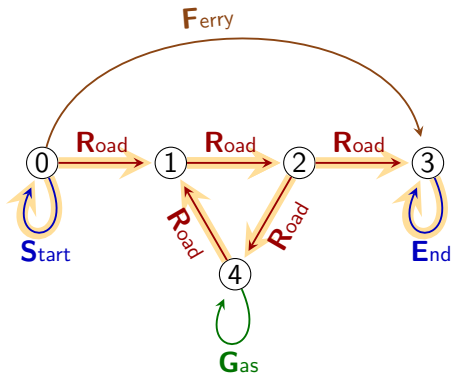
- The ferry
- The direct road



$$Q_1 = S(R+F)^*E$$

Q_1 matches...

- The ferry
- The direct road
- Roads with laps in the circuit



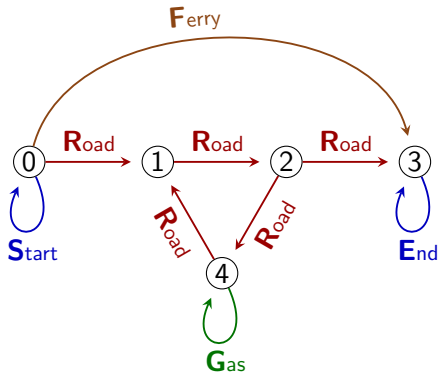
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Q_1 matches...

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- The direct road
- Roads with laps in the circuit

$$Q_2 = S(R+F)^*G(R+F)^*E$$

Q_2 matches...



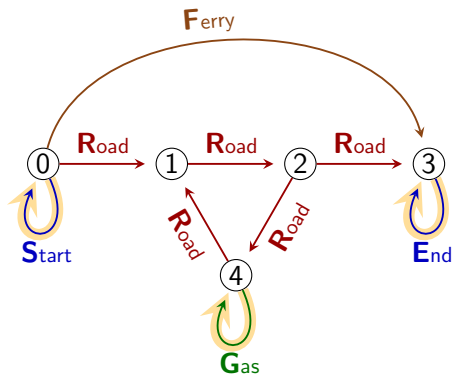
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$$Q_2 = S(R+F)^*G(R+F)^*E$$

Q_2 matches...



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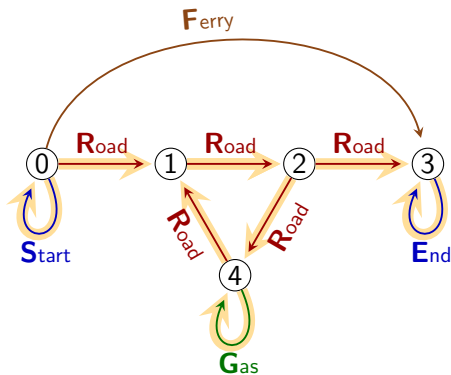
Q_1 matches...

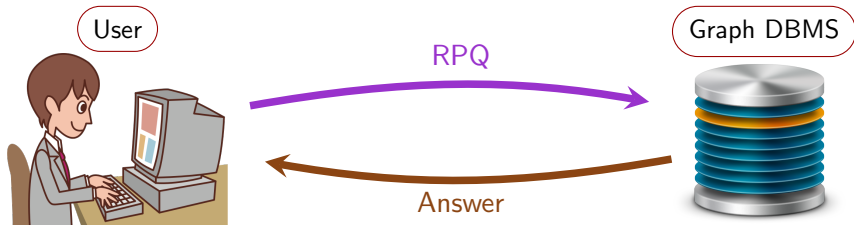
- The ferry
- The direct road
- Roads with laps in the circuit

$$Q_2 = S(R+F)^*G(R+F)^*E$$

Q_2 matches...

- Roads with laps in the circuit



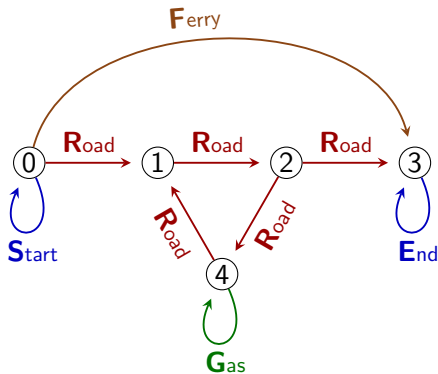


! Infinitely many matches but finite answer !

Main theoretical semantics [Angles et al. 2017] (used in SparQL)

Definition

- Returns the endpoints of matches



Main theoretical semantics [Angles et al. 2017] (used in SparQL)

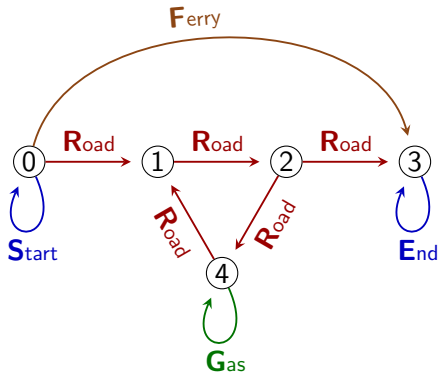
Definition

- Returns the endpoints of matches

$$Q_1 = S (R+F)^* E$$

$$Q_2 = S (R+F)^* G (R+F)^* E$$

- All matches are of the form:
 $0 \rightarrow \dots \rightarrow 3$
 $\Rightarrow Q_1$ and Q_2 return $\{(0, 3)\}$



Pros and cons

Pros

- Efficient algorithms
- Well grounded theory

Pros and cons

Pros

- Efficient algorithms
- Well grounded theory

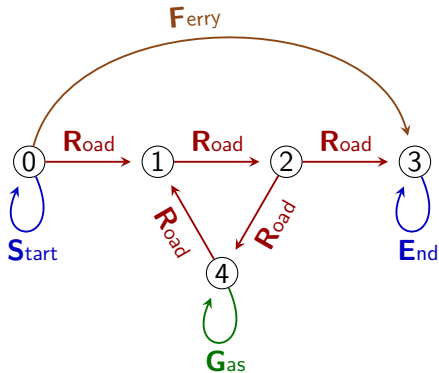
Cons

- Very limited information in the answer
 - User: *"I want to go from Paris to Lyon by car"*
 - Database: *"Yes you can"*

Used in PGQL (Oracle), in GSQL (TigerGraph), in G-core [Angles et al. 2018]

Definition

- Return the walk with the least number of edges



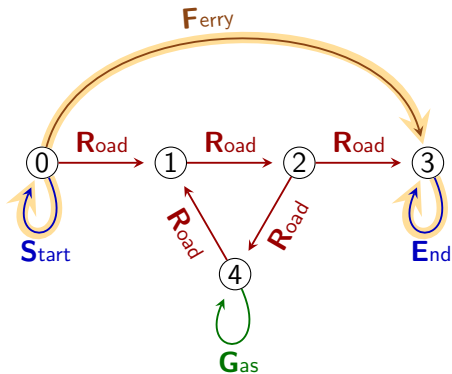
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Definition

- Return the walk with the least number of edges

$$Q_1 = S (R + F)^* E$$

- Q_1 returns 1 walk
 - the ferry
- Walks taking the road have more edges



Used in PGQL (Oracle), in GSQL (TigerGraph), in G-core [Angles et al. 2018]

Definition

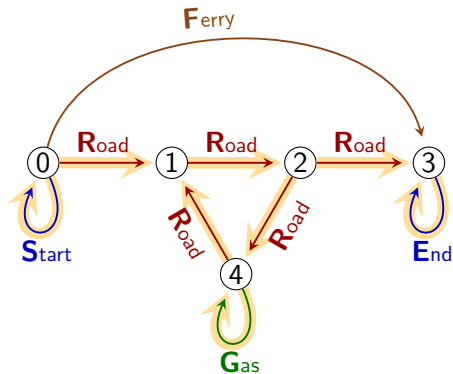
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 - the ferry
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$$Q_2 = S (R+F)^* G (R+F)^* E$$

- Q_1 returns 1 walk
 - the one-lap road



Pros and cons

Pros

- Returns walks
- Efficient algorithms
- Horizontal post-processing
 - Horizontal = along the walk
 - *“Is there a gas station on the way?”*
 - *“What is the length of the walk?”*

Pros and cons

Pros

- Returns walks
- Efficient algorithms
- Horizontal post-processing
 - Horizontal = along the walk
 - *“Is there a gas station on the way?”*
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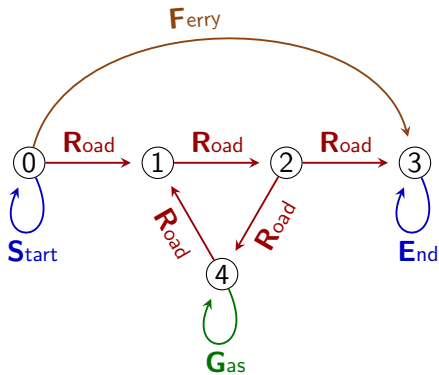
Cons

- No vertical post-processing
 - Vertical = accross the walks with the same endpoints
 - *“What is the shortest walk in time?”*
 - *“What is the connectedness level?”*
- No coverage of the space of matches

Used in Cypher (Neo4j) [Francis et al. 2018] [Green et al. 2019]

Definition

- Return walks
- Forbid to repeat edges



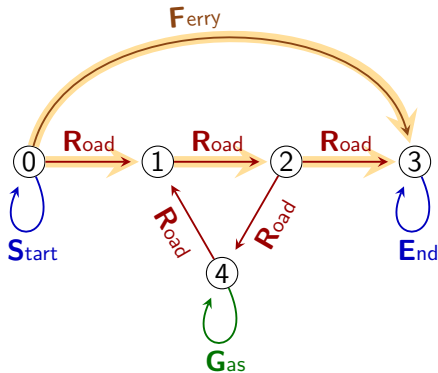
Used in Cypher (Neo4j) [Francis et al. 2018] [Green et al. 2019]

Definition

- Return walks
- Forbid to repeat edges

$$Q_1 = S (R + F)^* E$$

- Q_1 returns 2 walks
 - the ferry
 - the straight road



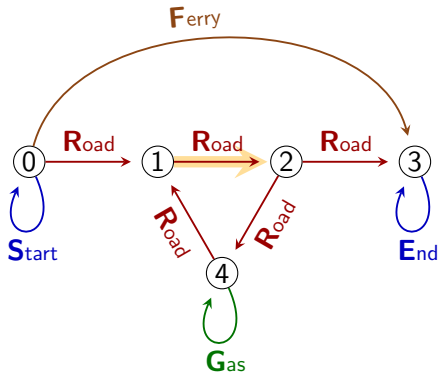
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Definition

- Return walks
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$$Q_1 = S (R + F)^* E$$

- Q_1 returns 2 walks
 - the ferry
 - the straight road
- Walks with circuit laps repeat the middle edge



Used in Cypher (Neo4j) [Francis et al. 2018] [Green et al. 2019]

Definition

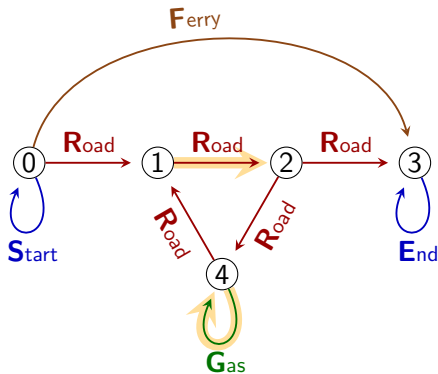
- Return walks
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$$Q_1 = S (R + F)^* E$$

- Q_1 returns 2 walks
 - the ferry
 - the straight road
- Walks with circuit laps repeat the middle edge

$$Q_2 = S (R + F)^* G (R + F)^* E$$

- Q_2 returns nothing



Pros and cons

Pros

- Returns walks
- Counting matches is possible
- Horizontal and vertical post-processing
- Some coverage of the space of matches

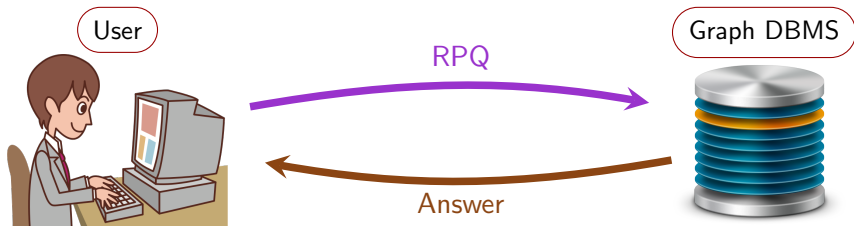
Pros and cons

Pros

- Returns walks
- Counting matches is possible
- Horizontal and vertical post-processing
- Some coverage of the space of matches

Cons

- All problems are computationally hard [Martens et al. 2020]
 - Counting, enumeration, existence
 - Checking whether Q_2 returns anything → Already NP-hard
- Part of the space of matches might be uncovered



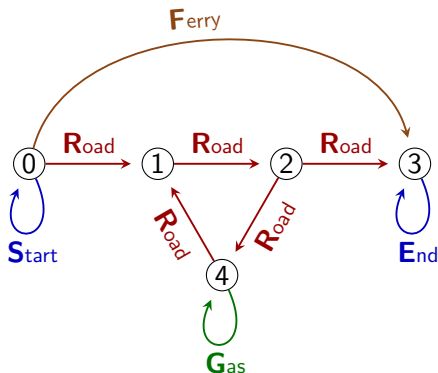
⚠ Infinitely many matches but finite answer ⚠

- Several way to ensure finiteness
 - Homomorphism → Filters out most information
 - Shortest-walk → Bad coverage of the space of matches
 - Trail → Computationally hard
 - Other variants have similar issues.
- No solution is clearly superior

New theoretical compromise [David-Francis-Marsault 202?]

Definition

- Returns walks
- Each edge may match each atom only once

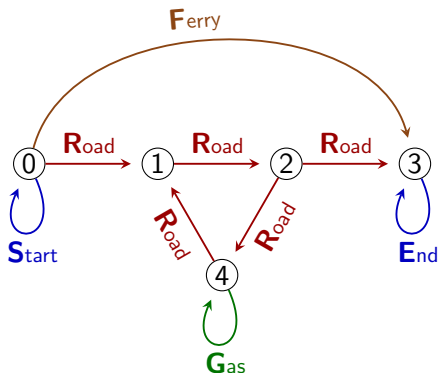


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$$Q_2 = S (R+F)^* G (R+F)^* E$$



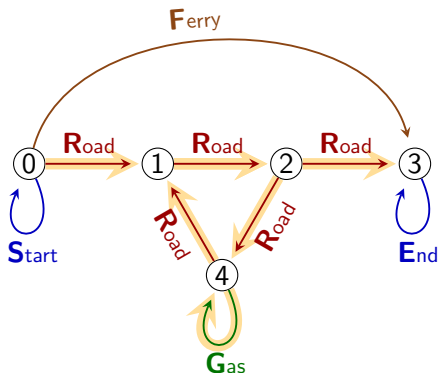
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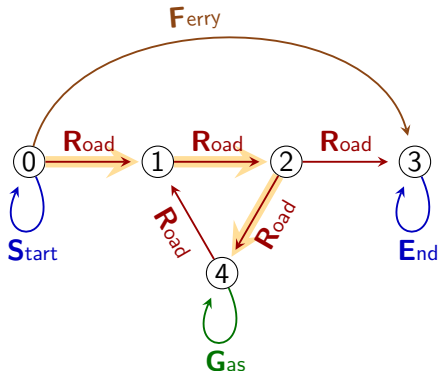
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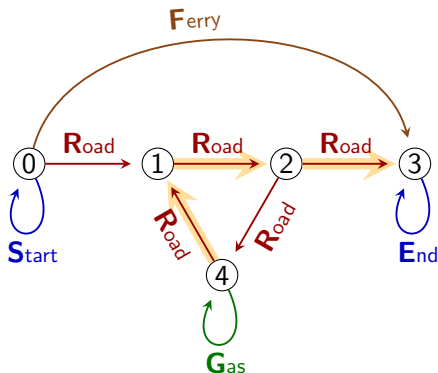
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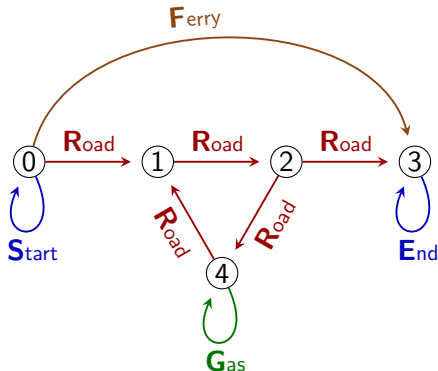
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 - Before **G** → use the left **R**
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- > 1 circuit lap ⇒ some edge use the same atom twice



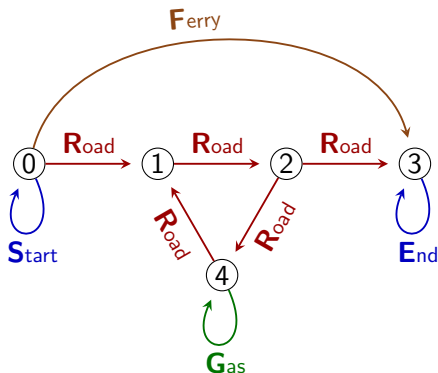
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 - Before $\mathbf{G} \rightarrow$ use the left \mathbf{R}
 - After $\mathbf{G} \rightarrow$ use the right \mathbf{R}
- > 1 circuit lap \Rightarrow some edge use the same atom twice



$$Q_1 = \mathbf{S} (\mathbf{R} + \mathbf{F})^* \mathbf{E}$$

- Returns the ferry and the straight road

Pros and cons

Pros

- Returns walks
- Horizontal and vertical post-processing
- "Reasonable" coverage of the space of matches
- Counting results is possible
- Emptiness and Enumeration are efficient

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- Counting results is computationally hard
- Answer depends on the way the query is written
 - \mathbf{R}^* allows no lap in the circuit
 - $(\mathbf{R} + \mathbf{R})^*$ allows 1 lap in the circuit

Pros and cons

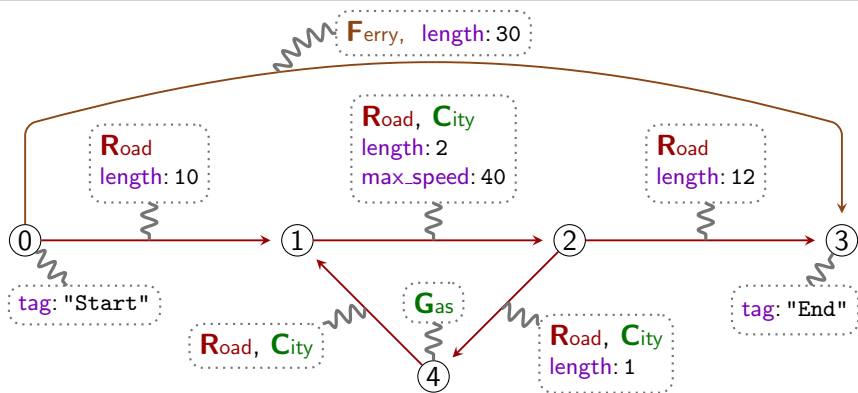
Pros

- Returns walks
- Horizontal and vertical post-processing
- "Reasonable" coverage of the space of matches
- Counting results is possible
- Emptiness and Enumeration are efficient
- Gives some expressivity to the user

Cons

- Counting results is computationally hard
- Answer depends on the way the query is written
 - \mathbf{R}^* allows no lap in the circuit
 - $(\mathbf{R} + \mathbf{R})^*$ allows 1 lap in the circuit

Property graphs and real query languages



Vertices and edges may bear:

- zero or more labels
- zero or more properties

- Property = key-value pair
- Key = string
- Value = bool, int, str, ...

- Trail semantics

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- Cypher is graph-to-tables
- Chaining of clauses

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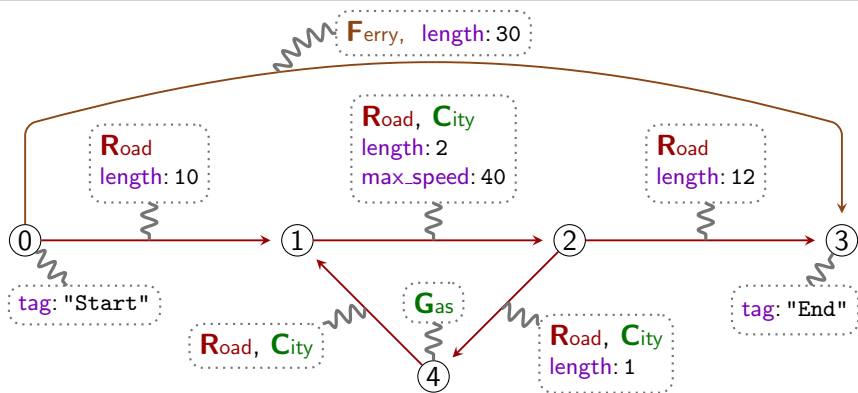
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Cypher queries for Q_1 and Q_2

```
MATCH ({tag:"Start"})-[:Road|Ferry*]->({tag:"End"})
```

```
MATCH ({tag:"Start"})-[:Road|Ferry*]->
      (:Gas)-[:Road|Ferry*]->({tag:"End"})
```

Cypher returns a table...



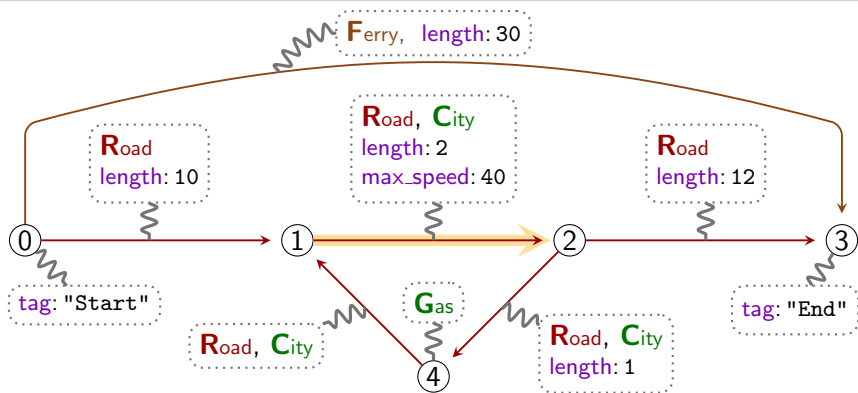
Query

```
MATCH (s)-[:City]->(t)
```

Result

s	t
1	2
2	4
4	1

Cypher returns a table...



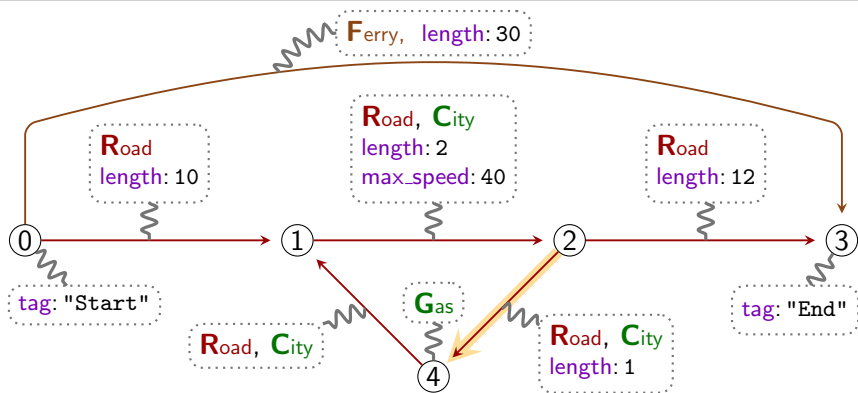
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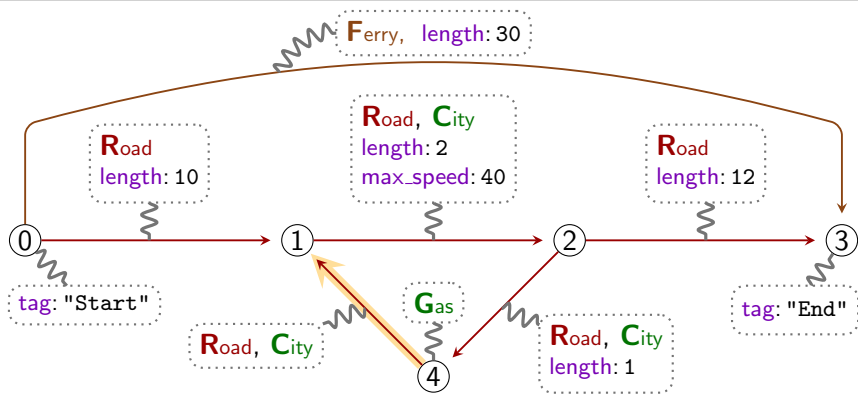
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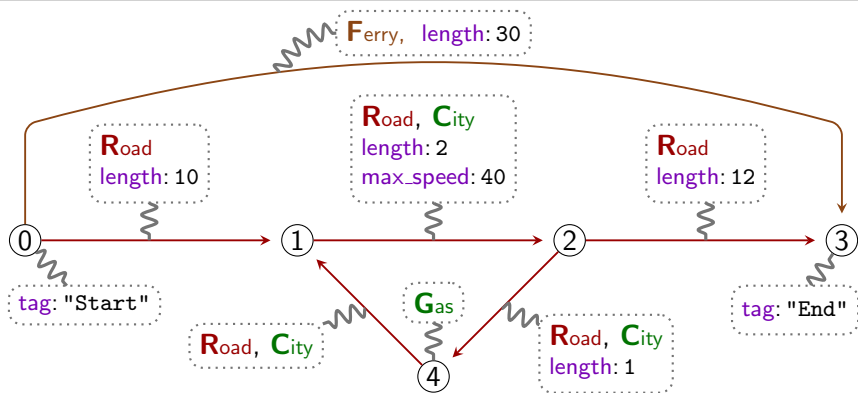
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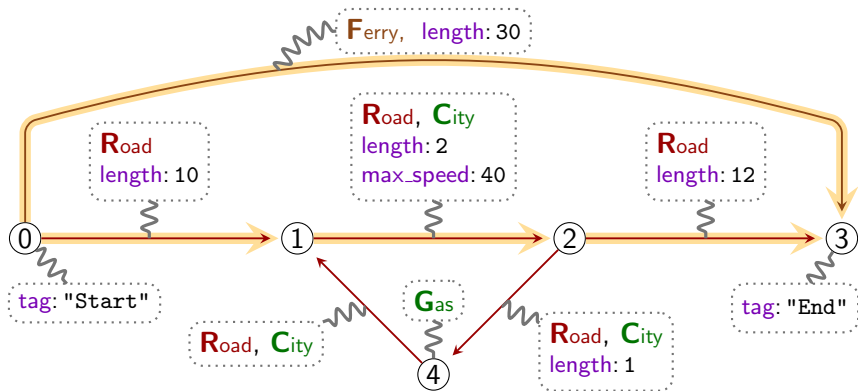
MATCH

```
(s {tag:"Start"})  
  -[:Road|Ferry*]->  
    (t {tag:"End"})
```

Result

s	t
0	3
0	3

Cypher returns a table... but computes walks



Query

MATCH

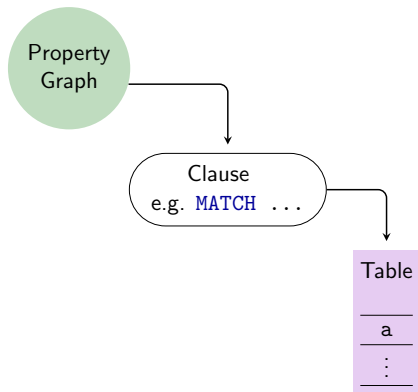
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-[:Road|Ferry*]->  
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```

Result

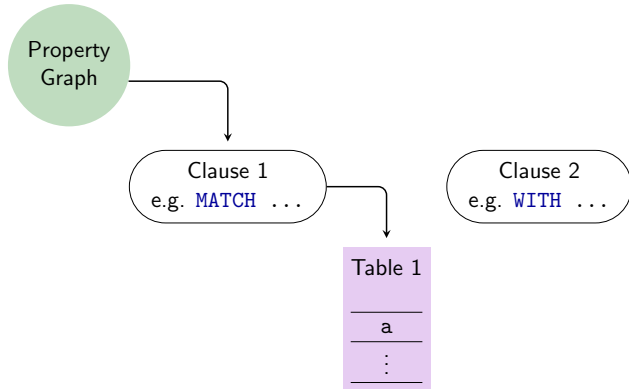
s	t	
0	3	← The ferry
0	3	← The direct road

- **ORDER BY:** orders row
- **WHERE:** filters row
- **WITH** or **RETURN:**
 - adds/renames columns
 - horizontal aggregation (e.g. with keyword reduce)
 - vertical aggregation (e.g. with keyword count, max)
- **CREATE/DELETE/SET:** updates the property graph

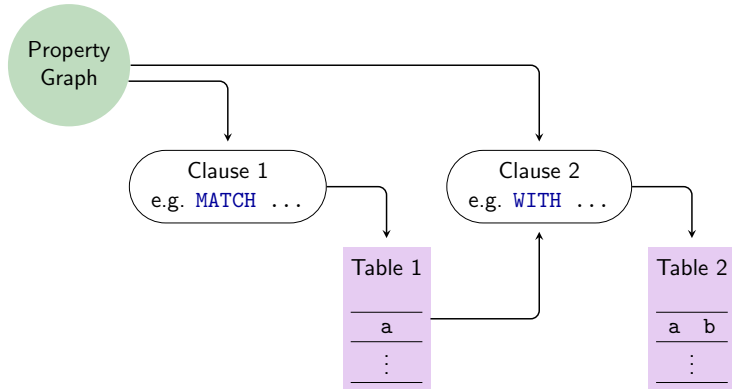
A Cypher query actually chain clauses



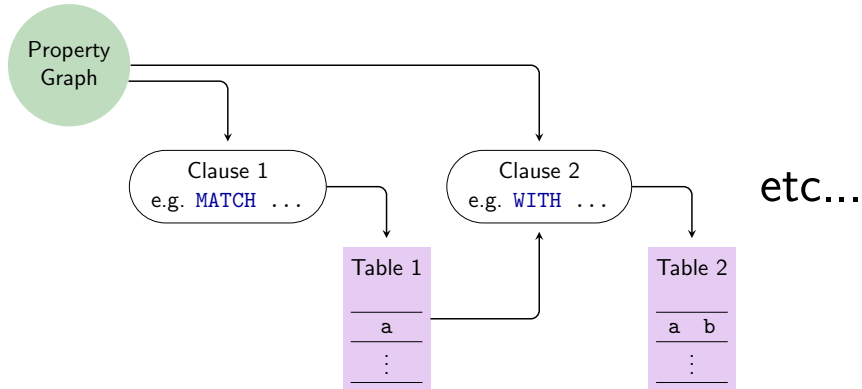
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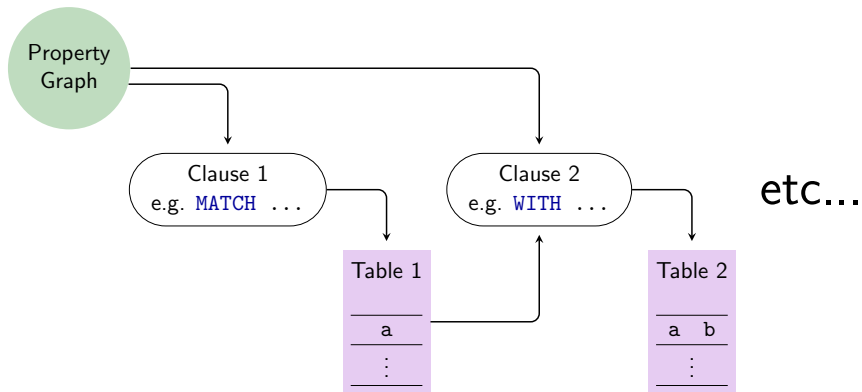


A Cypher query actually chain clauses



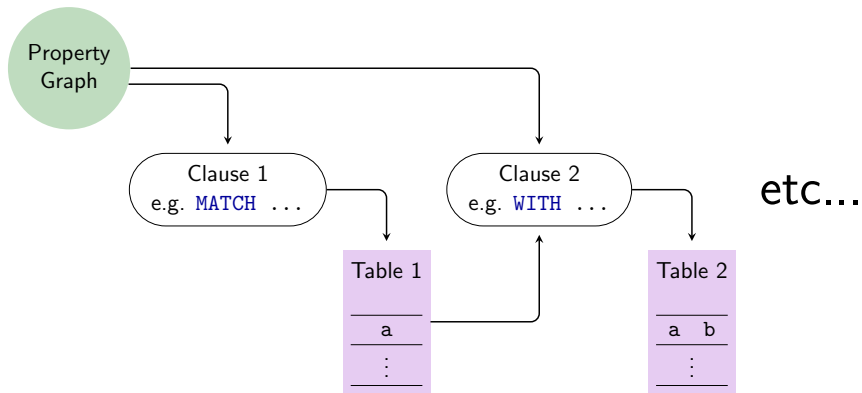
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Example

- Clause 1 makes some pattern matching
- Clause 2 aggregates over the result of Clause 1



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⇒ Trail semantics (rich post-processing at the cost of efficiency)

Features inherited from Cypher

- ASCII-art syntax
- Graph-to-tables
- Chaining of clauses
- No nested Kleene stars

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- Undirected edges
- Query multiple database at the same time
- Subqueries

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New features

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- Deduplication based on "binding path"

An RPQ may have infinitely many matches

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GQL does not choose

- Trail semantics → keyword **TRAIL**
- Shortest-walk semantics → keyword **SHORTEST**
- Syntax restriction → keyword **WALK**
- Others

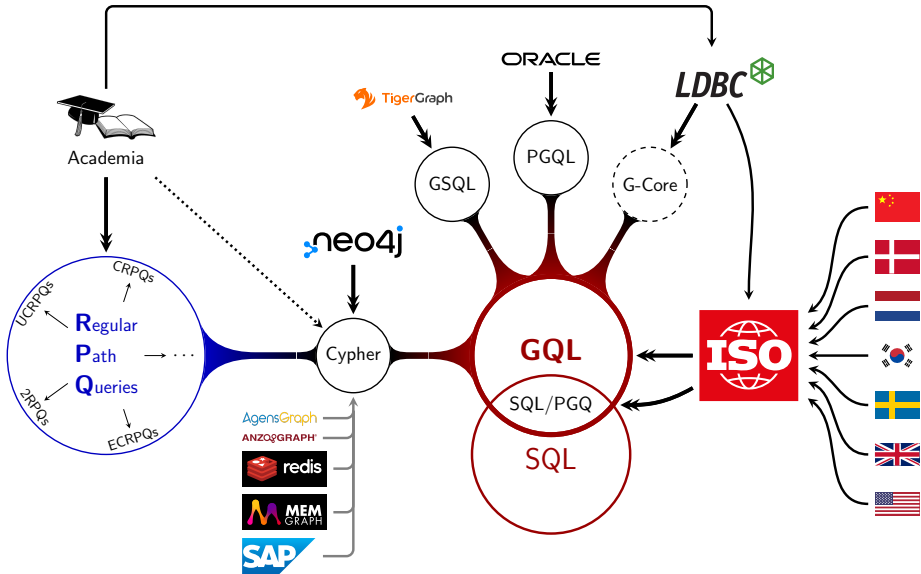
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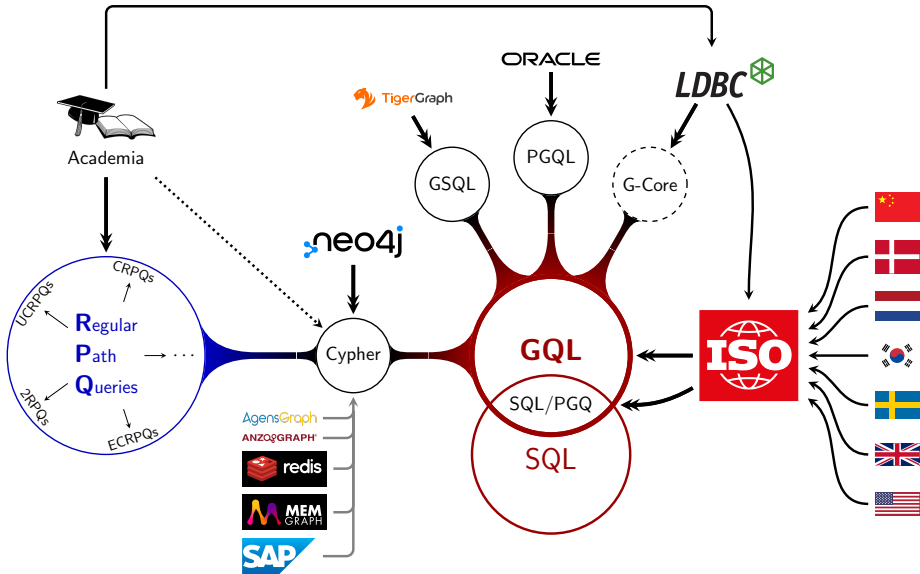
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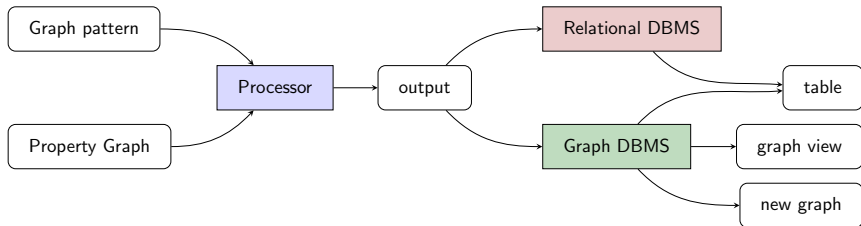
→ Could we add run-based semantics in GQL 2.0?

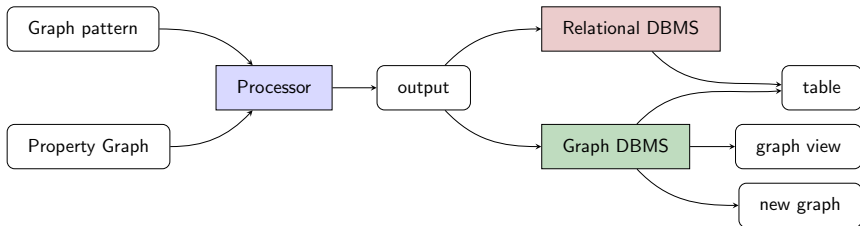




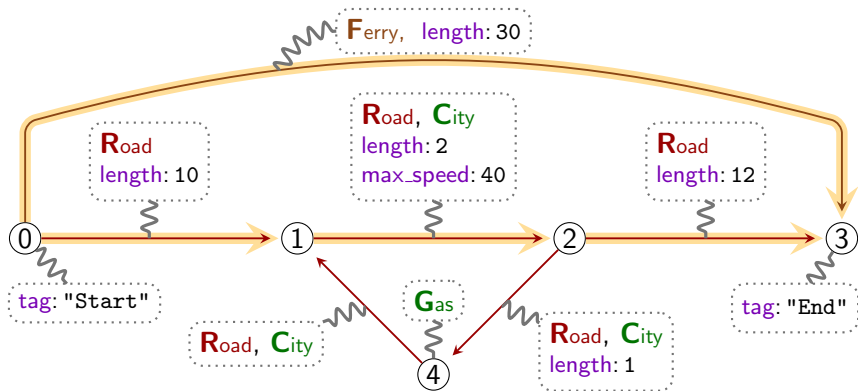
Thank you for your attention!

- Introduction
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- Foundation of querying graph databases: RPQs
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 - GQL 28



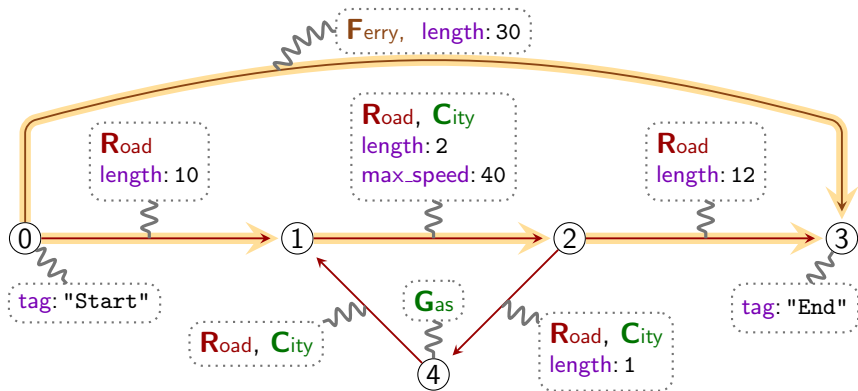


Output of GQL: set of path bindings
Path binding = walk annotated with variables



```
MATCH TRAIL (a WHERE a.tag="Start")
  [ -[r:Road]-> | -[c:City]-> ]* (b WHERE b.tag="End")
```

GQL path-bindings in one slide



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```

```
0 → 1 → 2 → 3  
a r  r  r b
```

```
0 → 1 → 2 → 3  
a r  c  r b
```